











Type-tagging with Forth

Formulating Type Tagged Parse Trees as
Forth Programs.

Dr Campbell Ritchie and Dr Bill Stoddart

Formal Methods and Programming Research Group
Teesside University, UK

A ::= A + T | A - T | T

A = A “+” T, A “-” T, T

A = A “+” T, A “-” T, T

T = T “*” F, T “/” F, F

F = N, I, “(” A “)”

• • •

$$P_A("a + t") = P_A(a) \ P_T(t) "+"$$

$$P_T("t * f") = P_T(t) \ P_F(f) " * "$$

$$P_F(" (a)") = P_A(a)$$

$$P_F("123") = P_N("123")$$

$$a + b * (c + d)$$

Becomes

$$a\ b\ c\ d\ +\ *\ +$$

$$\frac{a + b * (c + d)}{}$$

$$a \frac{b * (c + d)}{}$$

$$a b \frac{(c + d)}{*} +$$

$$a b c \frac{d}{*} +$$

$$a b c d + *$$

$1 + 2.5$

“ 1 ” “ INT ” “ 2.5 ” “ FLOAT ” + _

“ 1 FLOAT 2.5 F+ ”

Examples of types of Expressions:

No 1:

{ 1, 2, 3 } is a member of the
“Set of INTs” which is expressed in
postfix as
“ INT SET”

and the expression in postfix is
“ INT { 1 , 2 , 3 ,} INT SET”

Examples of types of Expressions:

No 2:

$1 \mapsto \text{“ foo”}$

is a member of INTs paired to STRINGs,
which can be expressed in postfix as
“ INT STRING PAIR”

and the expression in postfix is

“ 1 “ foo” \mapsto ”

$E = E \text{ ``} \mapsto \text{'' } E_0, E_0$

“ $x \mapsto y$ ”

“ x ” “foo” “ y ” “bar” $\mapsto_{_}$

“ $x \ y \ \mapsto$ ” “foo bar PAIR”

```
: $\mapsto$ _  
(  
    l-value:$ l-type:$ r-value:$ r-type:$ --  
    az1 = values catenated with  $\mapsto$   
    az2 = types catenated with PAIR  
)  
(: VALUE l-value VALUE l-type VALUE r-value VALUE r-type :)  
l-value r-value AZ^    "    $\mapsto$ " AZ^  
l-type r-type AZ^    "   PAIR" AZ^  
2LEAVE ;
```

The `,_` operation is used for creating lists for the contents of sets, etc.

It needs to check that the two types are *null* or the same type.

$$S = S_1 \text{ ``} \triangleleft \text{'' } S, S_1$$

$$\begin{array}{ccc} \mathbb{P}(T) & \triangleleft \mathbb{P}(T \times U) & \rightarrow \mathbb{P}(T \times U) \\ \text{Set of } T & \triangleleft \text{ Set of } T \text{ paired to } U & \end{array}$$

$$a \triangleleft b$$

$$\text{``} a \text{''} \text{ ``} \text{foo SET''} \text{ ``} b \text{''} \text{ ``} \text{foo bar PAIR SET''} \triangleleft \underline{\quad}$$

$$\text{``} a \ b \triangleleft \text{''} \text{ ``} \text{foo bar PAIR SET''}$$

“ foo bar PAIR SET”

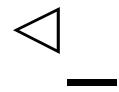
“ foo SET”

“ foo bar PAIR SET”

“ foo SET”

“ *a*” “ foo bar PAIR SET” “ *b*”

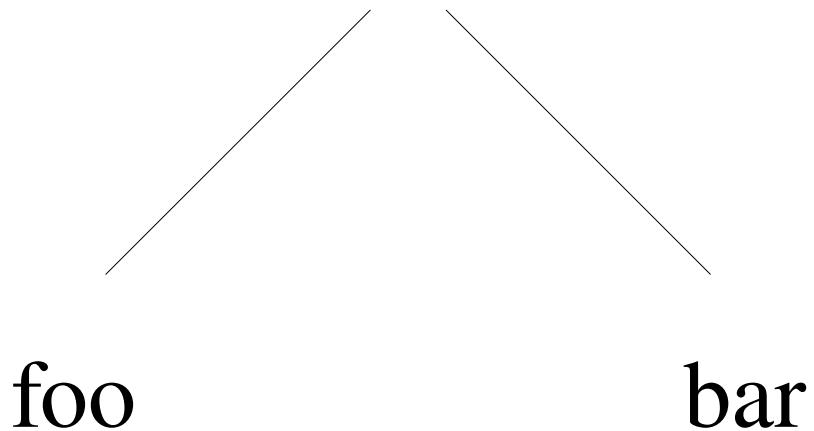
“ foo bar PAIR baz biz PAIR boz PAIR buz
PAIR SET PAIR SET”



SET



PAIR

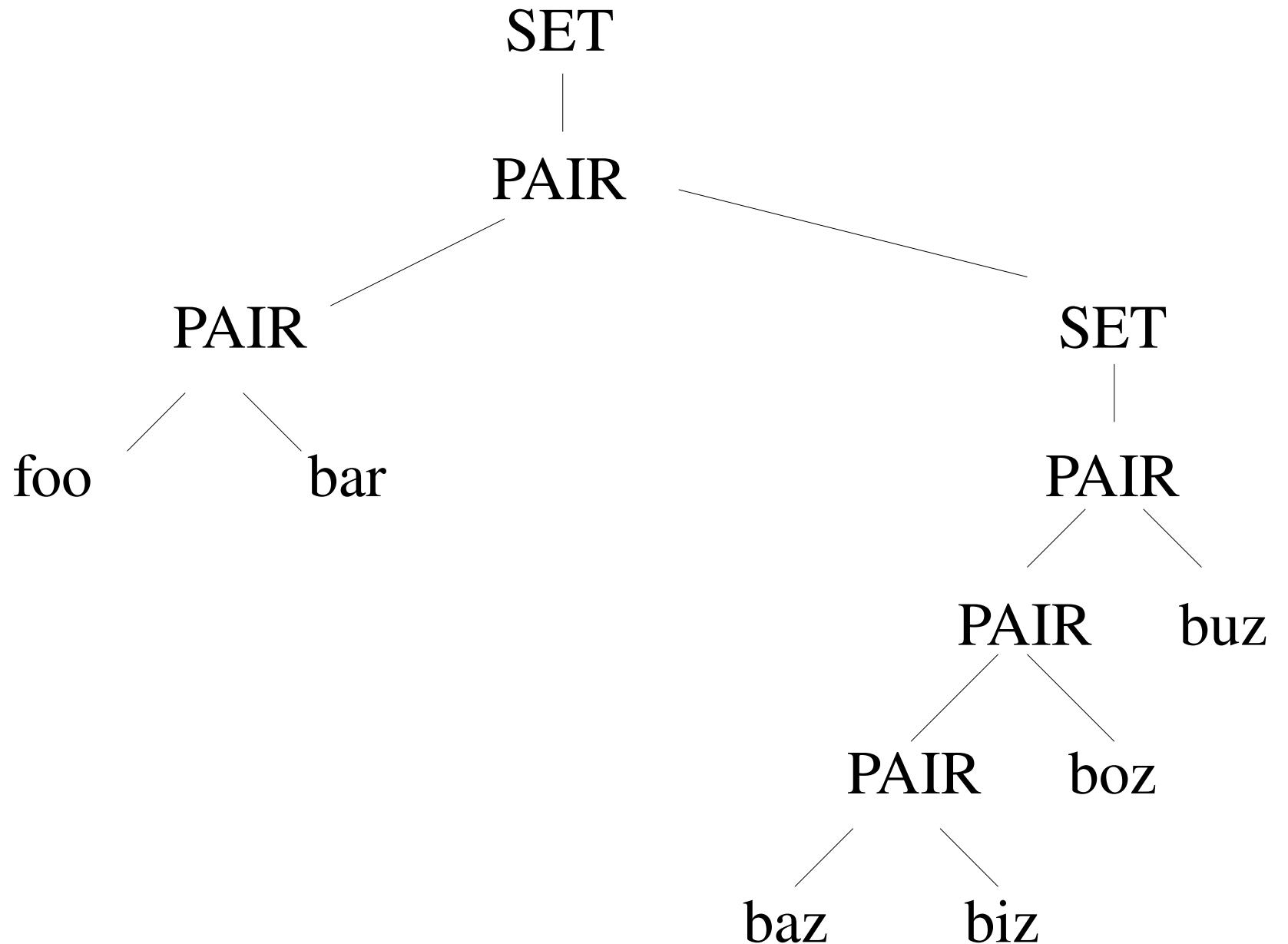


SET

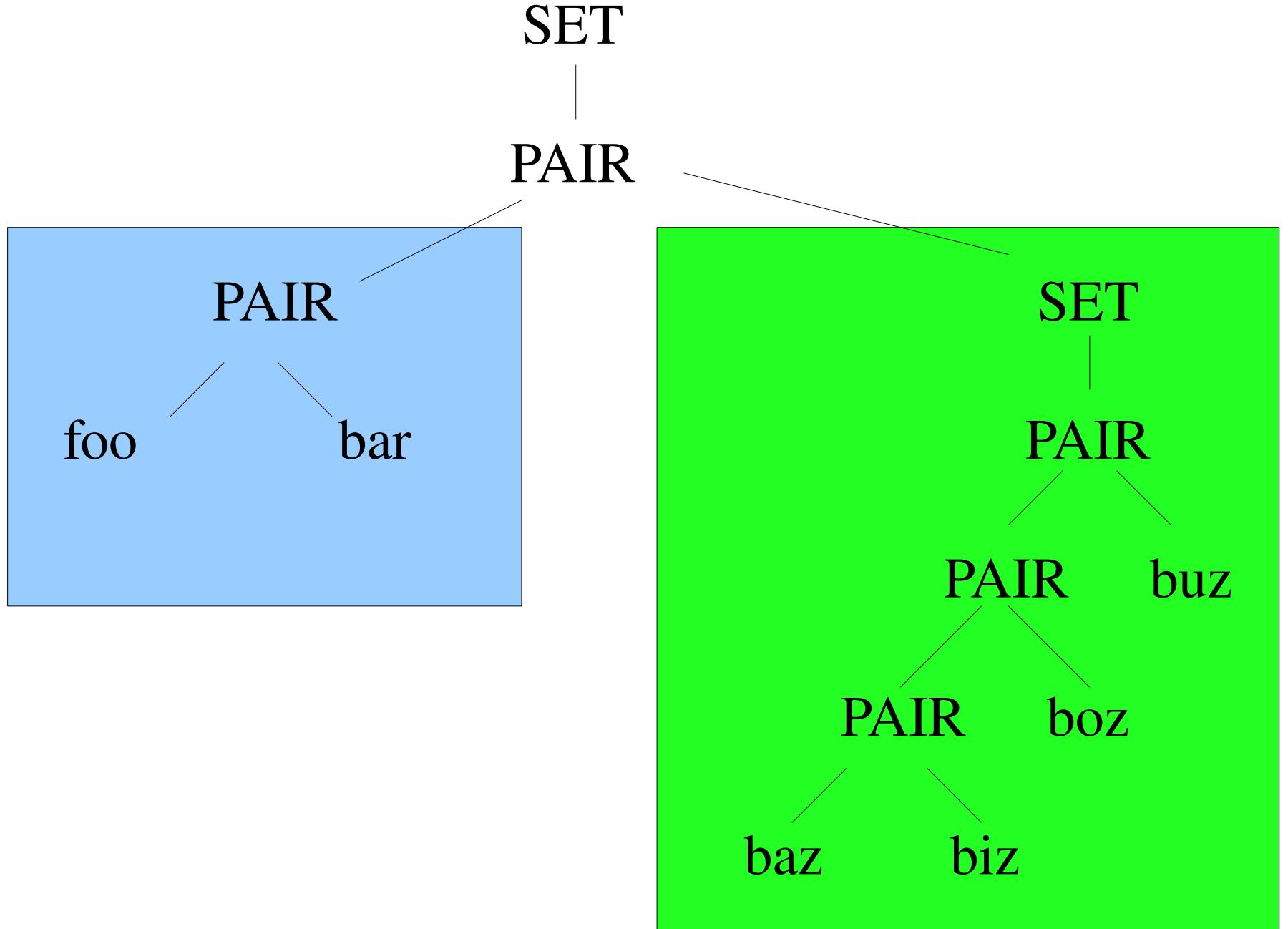
PAIR

foo

bar



“~~foo bar PAIR~~ baz biz PAIR boz PAIR buz
PAIR SET ~~PAIR SET~~”



“ *a*” “ foo bar PAIR SET” “ *b*”

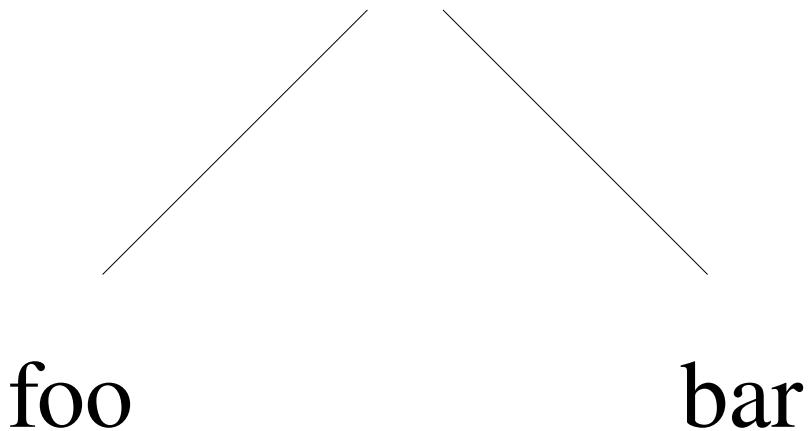
“ foo bar PAIR baz biz PAIR boz buz
PAIR SET PAIR SET”

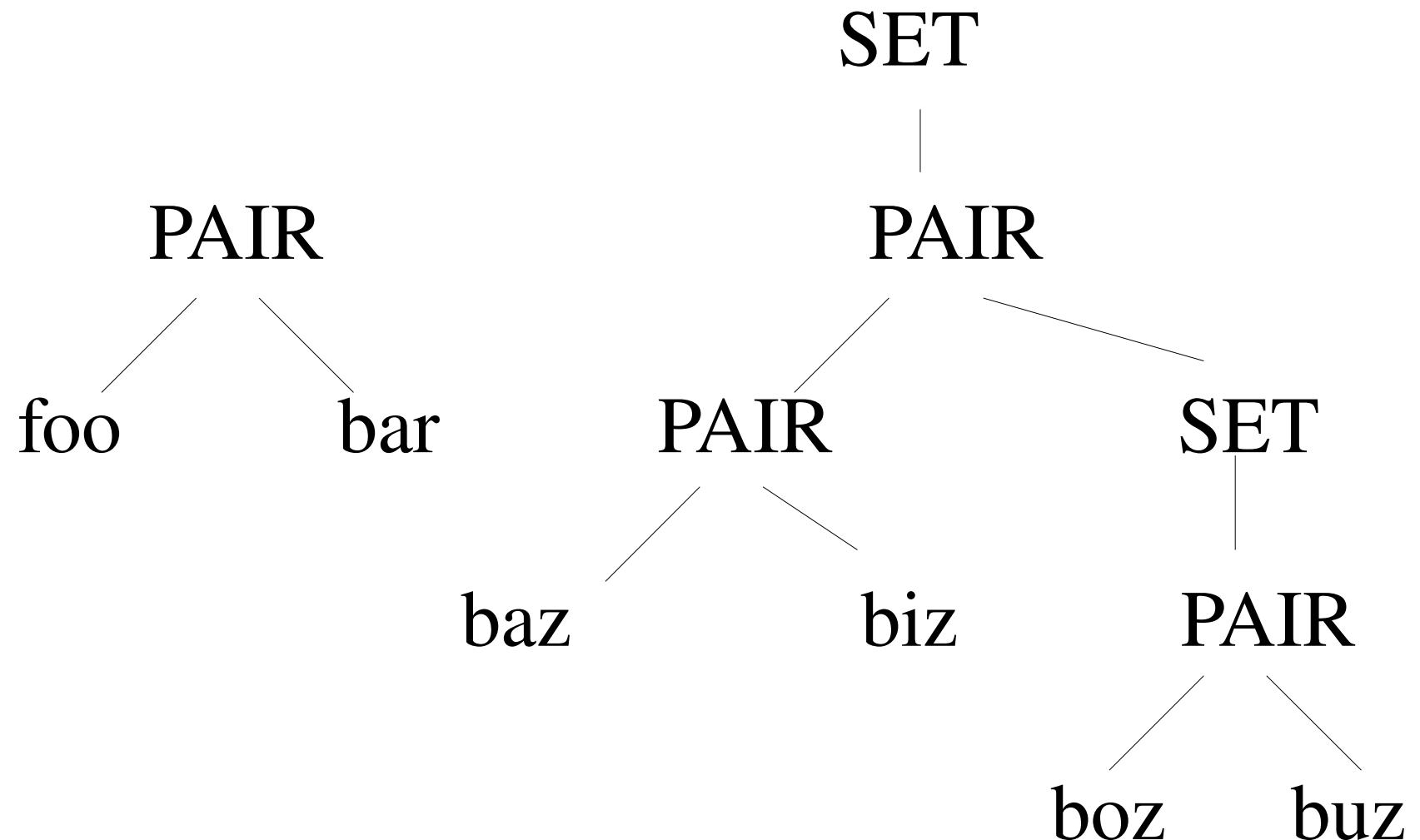
▷
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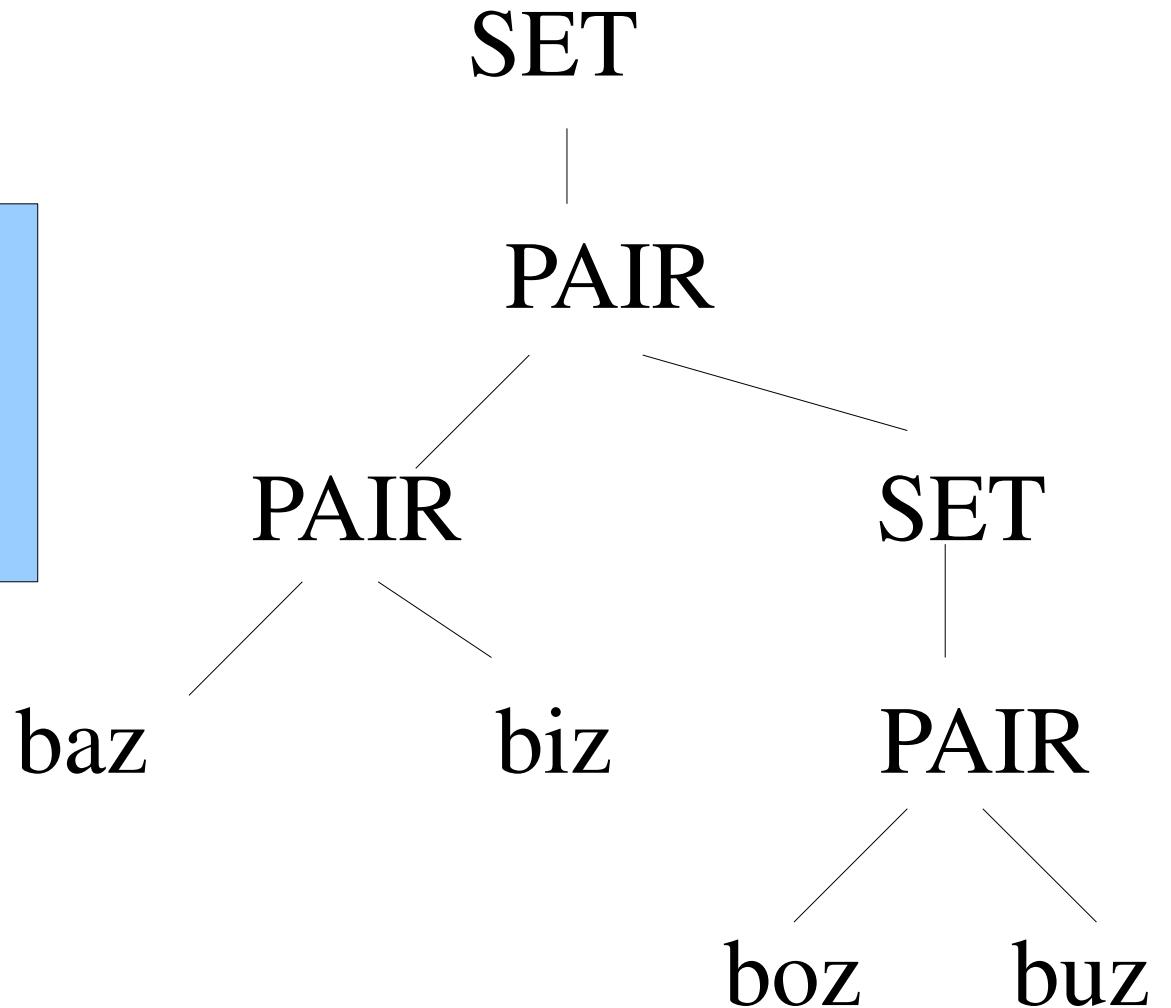
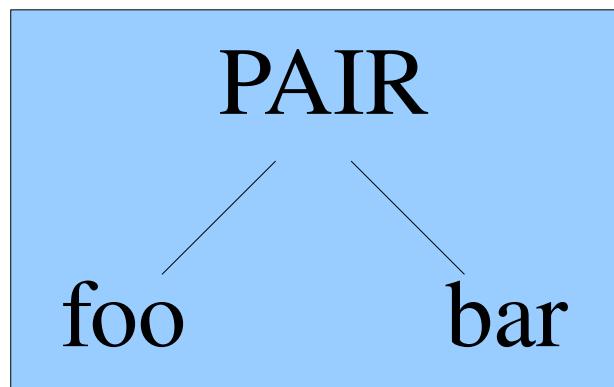
SET



PAIR







“ foo bar PAIR baz biz PAIR boz PAIR buz
PAIR SET ~~PAIR SET~~”

“ baz biz PAIR boz PAIR buz PAIR SET”

1 2 3-2=1 2 3-2=1 2 3-2=1 2-1=1

“ ~~foo bar PAIR~~ baz biz PAIR boz buz
PAIR SET ~~PAIR SET~~”

“ baz biz PAIR boz buz PAIR SET”

1 2 3-2=1 2 3 4-2=2 3-1=2

“ foo bar ~~PAIR~~ baz PAIR biz PAIR boz PAIR buz
bez PAIR SET ~~PAIR SET~~”

“ baz PAIR biz PAIR boz PAIR buz bezPAIR SET”

1 2-2=0 1 2-2=0 1 2-2=0 1 2 3-2=1 2-1=1

Conclusions

1: It is simple to add type Strings to FORTH expressions.

Conclusions

2: Analysing these types in a second pass of compilation allows the types of the result to be determined and correctness of the types passed to be verified.

Conclusions

3: We have a simple method for verifying that a postfix String represents a well-formed single expression.

